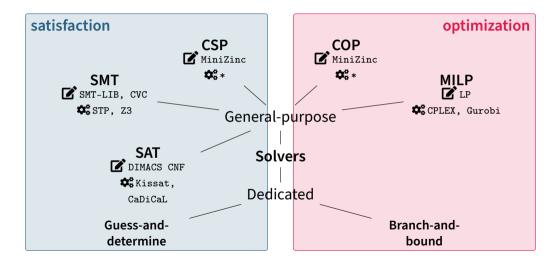


# Challenges in Tool-based Cryptanalysis

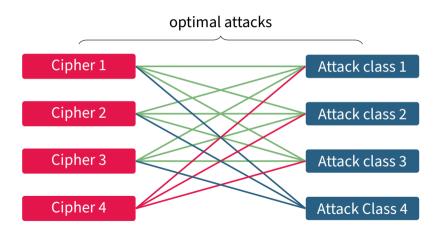
#### Maria Eichlseder

SKCAM 2025 Workshop on Symmetric-Key Cryptanalysis Automation and Modelling

### **Tools for Cryptanalysis**



#### The Ideal Framework



### Frameworks: The Challenges of Getting the Full Picture

- Cipher representation
  - What are the right abstraction layers? Which design paradigms to support?
  - Simplicity vs. versatility
- Precision of results
  - What to assume, what to analyze?
  - Efficiency vs. quality
- Scope of results
  - Building blocks or end-to-end attacks?
  - Optimality vs. feasibility

# Cipher Representation

Simplicity vs. Versatility

## Direction 1: The Right Abstraction Level

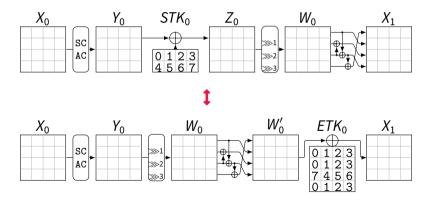
- Construction: Key-alternating SPN/LS, Feistel, Reflection, Alignment, ...
- Operations: S-box, MDS matrix, modular addition, permutation,  $\mathbb{F}_p, \dots$
- Gates and local operations:  $\oplus$ ,  $\odot$ , ...

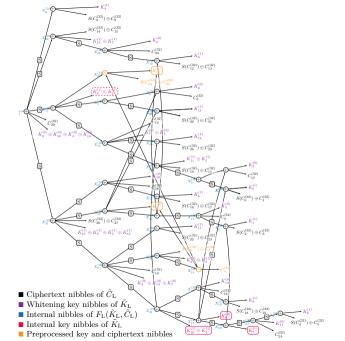
#### Decomposition into smaller functions:

- Necessary for efficient modelling
- Loses information (invertibility, MDS property, ...), introduces inaccuracies

# **Direction 2: Rewiring the Cipher**

- Restructure the cipher circuit representation for optimizations
- **Examples:** Equivalent subkeys, Round boundary,  $f(C \oplus K)$  for FFT, ...



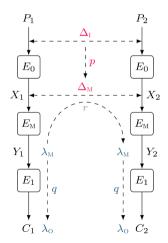


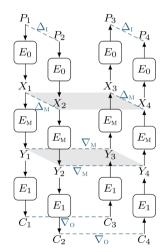
# Precision of Results

Efficiency vs. Quality

#### **Direction 3: Combined Trails**

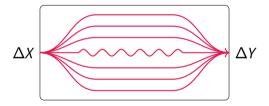
**Examples:** Differential-linear, Boomerang, Miss-in-the-middle, ...





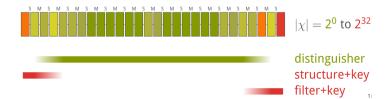
#### **Direction 4: Abundance of Trails**

- More precise results by analyzing all relevant trails
- Leads to 2-level optimization problems
- > Examples: Differential clustering, Linear hull, Quasidifferential trails



#### Direction 4: Abundance of Trails

- More precise results by analyzing all relevant trails
- Leads to 2-level optimization problems
- **Examples:** Differential clustering, Linear hull, Quasidifferential trails
- > Representing partial knowledge and sets of trails
- **Examples:** (Partially) truncated propagation, Guess-and-determine



#### Direction 4: Abundance of Trails

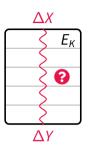
- More precise results by analyzing all relevant trails
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- **Examples:** Differential clustering, Linear hull, Quasidifferential trails
- Representing partial knowledge and sets of trails
- Examples: (Partially) truncated propagation, Guess-and-determine

$\nabla(z_j,z_j^*)$	$\nabla(z_j,z_j^*)$	$\nabla(z_j,z_j^*)$	$\nabla(z_j,z_j^*)$
0 = 000●	-=•○○●	3 = ○○●●	7 = ○●●●
u = 00●0	X = ○●●○	5 = ○●○●	B = •○••
n = ○●○○	#=0000	A = •○•○	$D = \bullet \bullet \circ \bullet$
1 = ●000	? = ●●●●	C = ••oo	E = ●●●○

#### Direction 5: Absence of Trails

Some distinguishers are based on the non-existence of a valid trails

- Differential > Impossible differentials
- Linear > Zero-correlation linear approximations
- Integral > Division/monomial trail; ZC-based integrals



However, models for full attacks need solution-based distinguisher models (or a quantified language like QSAT.)

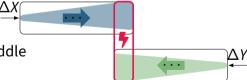
Proving absence of absence-based distinguishers?

#### Two Ways of Modelling Impossibility

- Unsatisfiability-based:
- > First specify distinguisher, then check
- > Precise, but potentially slow



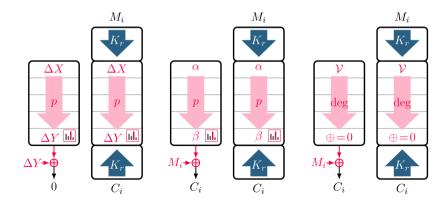
- Satisfiability-based:
- Find distinguisher that misses in the middle
- > Typically efficient, but less precise



# Optimizing Full Attacks

#### Block ciphers:

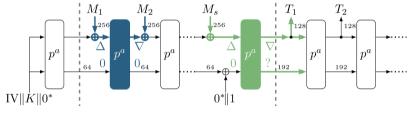
- Distinguisher (differential, linear, integral, combined trail?)
- Key recovery (algorithm...?)



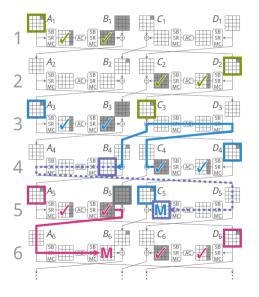
- Block ciphers:
  - Distinguisher (differential, linear, integral, combined trail?)
  - Key recovery (algorithm...?)
- Hash functions, compression functions:
  - ???
- Permutations:
  - ???
- Arithmetization-oriented primitives:
  - ???

We need to know

- which space to search
- what we can vary
- > what we can deduce



(a) Forgery (■) or MAC forgery (■) via differential cryptanalysis



- 1 Satisfy 1 and 2 fixed free
- 2 Satisfy 5 and 6 fixed match input
- Satisfy 3 and 4
  fixed
  match input
- ---> works because:
  - $\blacksquare$  fixed  $\rightarrow$   $\blacksquare$  known
  - $\blacksquare$  fixed  $\rightarrow$   $\blacksquare$  known

#### Optimizing Key-Recovery: Examples

Guess-and-determine attacks:

✓/> Autoguess [HE22a] ♠ a

tool

ahttps://github.com/hadipourh/autoguess

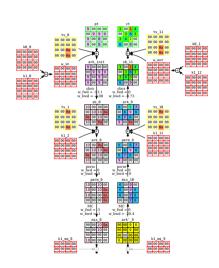
bhttps://extgit.isec.tugraz.at/castle/tool/keyrecoverytool

chttps://github.com/hadipourh/mpt

dhttps://github.com/hadipourh/zero

#### Optimizing Key-Recovery: Examples

- Guess-and-determine attacks:
  - Autoguess [HE22a] a
- Differential cryptanalysis:
  - keyrecoverytool [Nag22] \$\infty\$^b\$



ahttps://github.com/hadipourh/autoguess

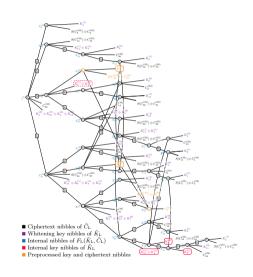
bhttps://extgit.isec.tugraz.at/castle/tool/keyrecoverytool

chttps://github.com/hadipourh/mpt

dhttps://github.com/hadipourh/zero

#### Optimizing Key-Recovery: Examples

- Guess-and-determine attacks:
   Autoguess [HE22a] a
- Integral cryptanalysis:
   Graph-based [HE22b]



ahttps://github.com/hadipourh/autoguess

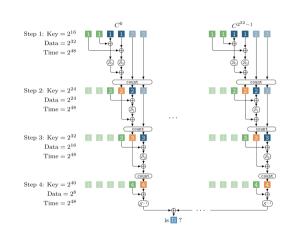
bhttps://extgit.isec.tugraz.at/castle/tool/keyrecoverytool

 $<sup>^{</sup>c}$ https://github.com/hadipourh/mpt

dhttps://github.com/hadipourh/zero

#### Optimizing Key-Recovery: Examples

- Integral cryptanalysis:
   Graph-based [HE22b] c
   AutoPSy [HSE23] c



ahttps://github.com/hadipourh/autoguess

bhttps://extgit.isec.tugraz.at/castle/tool/keyrecoverytool

chttps://github.com/hadipourh/mpt

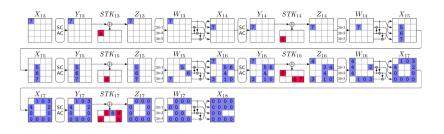
dhttps://github.com/hadipourh/zero

## **Direction 8: Optimizing Full Attacks**

- Key recovery has long been ignored
- Fewer choices to make for the attacker

...but ...

- Optimizations involve choices and tradeoffs
- Precise evaluation is tedious
- Optimal" distinguisher doesn't guarantee optimal attack



$\operatorname{Step}$	Guessed	$K\times D=\!\!\operatorname{Mem}$	$_{ m Time}$	Stored Texts
0	_	$2^0 \times 2^{40} = 2^{40}$		$Z_{17}[1,3,4,7]; X_{17}[8,11,12,13,15]; X_{16}[15]$
1	$STK_{17}[1]$	$2^4 \times 2^{36} = 2^{40}$	$2^{44-7.2}$	$Z_{17}[3,4,7]; X_{17}[8,11,12,15]; X_{16}[14,15]$
2	$STK_{17}[7]$	$2^8 \times 2^{32} = 2^{40}$	$2^{44-8.2}$	$Z_{17}[3,4]; X_{17}[8,12,15]; Z_{16}[6]; X_{16}[14,15]$
3	$STK_{17}[3]$	$2^{12} \times 2^{28} = 2^{40}$	$2^{44-7.2}$	$Z_{17}[4]; X_{17}[8,12]; Z_{16}[6]; X_{16}[12,14,15]$
4	$STK_{17}[4]$	$2^{16} \times 2^{28} = 2^{44}$	$2^{44-7.2}$	$Z_{16}[0,6,7]; X_{16}[10,12,14,15]$
5	$STK_{16}[6]$	$2^{20} \times 2^{20} = 2^{40}$	$2^{48-7.2}$	$Z_{16}[0,7]; X_{16}[12,15]; X_{15}[5]$
6	$STK_{16}[7]$	$2^{24} \times 2^{16} = 2^{40}$	$2^{44-7.2}$	$Z_{16}[0]; X_{16}[12]; X_{15}[5, 9]$
7	$STK_{16}[0]$	$2^{28} \times 2^4 = 2^{32}$		$X_{13}[0]$
$\Sigma$		$2^{44}$	$2^{41.32}$	

#### Why Optimizing Full Attacks is Challenging

- Preferably use a joint model for distinguisher and key recovery
  - ➤ Only works for satisfiability-based distinguishers
- ☐ Complexity formulas are often complicated
  - ➤ Mix of polynomial/exponential terms; simplified assumptions
- Multi-step processes lead to heavy models
- Very different types of key schedules
- Many different optimizations and strategies

#### Conclusion

- **Output** Cipher representation
  - Simplicity vs. versatility
- Precision of results
  - Efficiency vs. quality
- Scope of results
  - Optimality vs. feasibility

#### **Acknowledgements**





This research was funded in part by the Austrian Science Fund (FWF) SFB project SPyCoDe (doi:10.55776/F85) and the European Research Council (ERC) starting grant KEYLESS (#101165216).

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